

# Vertically and Horizontally Interacting Newsvendors

## **Abstract**

In this working paper general conditions for a global optimum of the production quantity of several vertically respectively horizontally interacting newsvendors are derived.

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## 1 Vertically interacting newsvendors

We propose a supply chain structure of  $N$  value stages, where every production site is responsible for one value stage. Horizontal interaction is not considered. We further assume a typical newsvendor situation with a normal distributed demand  $D \sim N(\mu, \sigma)$ , where  $\mu$  represents the mean and  $\sigma$  denotes the standard deviation. The short sales season allows only one order point, capacities are unrestricted, and decision-makers behave rational. We further consider one season and one product. In order to produce a certain production quantity on a value stage, the same quantity of the upstream value stage is required. Hence, defectives are not considered. As all relevant production processes have to be finished before the sales season, lead times can be neglected. We further assume that a production process started at site 1 have to be finished at site  $N$ . Each site's production costs are given by  $c_1$  to  $c_N$ . The revenue  $r$  is paid by end customers for one unit of the product. The salvage value  $v_0$  of a product is the value of a product after the season.

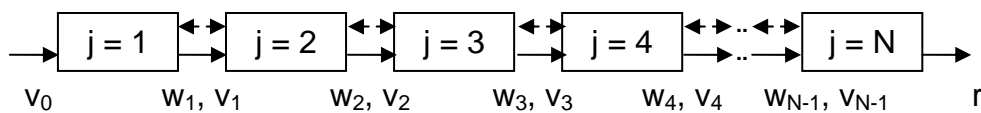


Figure 1: Decision situation for vertical interaction

The supply chain can be coordinated by the transfer prices  $w_1$  to  $w_{N-1}$ , where  $w_1$  represents the price, which site 1 achieves from site 2 for one unit. The buyback prices  $v_1$  to  $v_{N-1}$  are paid by the supplying site for every unit, which has not been sold during sales season.

*Theorem 1:*

If  $w_1, \dots, w_{N-1}$  and  $v_1, \dots, v_{N-1}$  are determined in a way that  $CR_j = CR_{j-1}$  holds for every  $j$  between 2 and  $N$ , then  $CR_j = CR_g$  is true for every  $j$  between 1 and  $N$ , where

$$CR_1 = \frac{w_1 - c_1}{v_1 - v_0} \quad (1)$$

and

$$CR_j = \frac{w_j - w_{j-1} - c_j}{v_j - v_{j-1}} \quad \text{for } 2 \leq j \leq N-1 \quad (2)$$

and

$$CR_N = \frac{r - w_{N-1} - c_N}{r - v_{N-1}} \quad (3)$$

and

$$CR_g = \frac{r - \sum_{k=1}^N c_k}{r - v_0}. \quad (4)$$

*Proof:*

First it is shown by induction that the statement

$$(CR_m = CR_{m-1} \quad \forall 2 \leq m \leq j) \Leftrightarrow v_j = \frac{v_1 - v_0}{w_1 - c_1} \left( w_j - w_1 - \sum_{k=2}^j c_k \right) + v_1 \quad (5)$$

is true for every  $j$  between 2 and  $N-1$ .

Basis: Showing that the statement (5) holds for  $j = 2$  and  $j = 3$ .

$j = 2$ :

$$\begin{aligned} CR_2 &= CR_1 \\ \Leftrightarrow \frac{w_2 - w_1 - c_2}{v_2 - v_1} &= \frac{w_1 - c_1}{v_1 - v_0} \\ \Leftrightarrow v_2 &= \frac{v_1 - v_0}{w_1 - c_1} (w_2 - w_1 - c_2) + v_1 \end{aligned} \quad (6)$$

$j = 3$ :

$$\begin{aligned}
CR_3 &= CR_2 \\
\Leftrightarrow v_3 &= \frac{v_2 - v_1}{w_2 - w_1 - c_2} (w_3 - w_2 - c_3) + v_2
\end{aligned} \tag{7}$$

Replacing  $v_2$  with the expression from (6) gives

$$\begin{aligned}
v_3 &= \frac{(v_1 - v_0)(w_2 - w_1 - c_2)}{(w_1 - c_1)(w_2 - w_1 - c_2)} (w_3 - w_2 - c_3) + \frac{v_1 - v_0}{w_1 - c_1} (w_2 - w_1 - c_2) + v_1 = \\
&= \frac{v_1 - v_0}{w_1 - c_1} (w_3 - w_2 - c_3 + w_2 - w_1 - c_2) + v_1 = \\
&= \frac{v_1 - v_0}{w_1 - c_1} (w_3 - w_1 - c_3 - c_2) + v_1
\end{aligned} \tag{8}$$

Inductive step:

Induction hypothesis: Statement (5) holds for  $j-1$  and for  $j$ .

Showing that Statement (5) then also holds for  $j+1$ :

$$\begin{aligned}
CR_{j+1} &= CR_j \\
\Leftrightarrow v_{j+1} &= \frac{v_j - v_{j-1}}{w_j - w_{j-1} - c_j} (w_{j+1} - w_j - c_{j+1}) + v_j
\end{aligned} \tag{9}$$

Using the induction hypothesis  $v_{j+1}$  can be transformed to

$$\begin{aligned}
v_{j+1} &= \frac{(v_1 - v_0)(w_j - w_{j-1} - c_j)}{(w_j - w_{j-1} - c_j)(w_1 - c_1)} (w_{j+1} - w_j - c_{j+1}) + \frac{v_1 - v_0}{w_1 - c_1} \left( w_j - w_1 - \sum_{k=2}^j c_k \right) + v_1 = \\
&= \frac{v_1 - v_0}{w_1 - c_1} \left( w_{j+1} - w_j - c_{j+1} + w_j - w_1 - \sum_{k=2}^j c_k \right) + v_1 = \\
&= \frac{v_1 - v_0}{w_1 - c_1} \left( w_{j+1} - w_1 - \sum_{k=2}^{j+1} c_k \right) + v_1
\end{aligned} \tag{10}$$

This way, it is proven that statement (5) holds for every  $j$  between 2 and  $N-1$ . We can

now calculate an expression for  $v_1$  from  $CR_N = CR_{N-1}$ :

$$\begin{aligned}
CR_N &= CR_{N-1} \\
\Leftrightarrow \frac{r - w_{N-1} - c_N}{r - v_{N-1}} &= \frac{w_{N-1} - w_{N-2} - c_{N-1}}{v_{N-1} - v_{N-2}} \\
\Leftrightarrow (r - w_{N-1} - c_N)(v_{N-1} - v_{N-2}) &= (w_{N-1} - w_{N-2} - c_{N-1})(r - v_{N-1})
\end{aligned} \tag{11}$$

Replacing  $v_{N-1}$  and  $v_{N-2}$  with the expression from (5):

$$\begin{aligned}
&\Leftrightarrow (r - w_{N-1} - c_N) \frac{v_1 - v_0}{w_1 - c_1} (w_{N-1} - w_{N-2} - c_{N-1}) = \\
&= (w_{N-1} - w_{N-2} - c_{N-1}) \left[ r - \frac{v_1 - v_0}{w_1 - c_1} \left( w_{N-1} - w_1 - \sum_{k=2}^{N-1} c_k \right) - v_1 \right] \\
&\Leftrightarrow (r - w_{N-1} - c_N) \frac{v_1 - v_0}{w_1 - c_1} = r - \frac{v_1 - v_0}{w_1 - c_1} \left( w_{N-1} - w_1 - \sum_{k=2}^{N-1} c_k \right) - v_1 \\
&\Leftrightarrow r(w_1 - c_1) = \\
&= (v_1 - v_0) \left( w_{N-1} - w_1 - \sum_{k=2}^{N-1} c_k \right) + v_1(w_1 - c_1) + (r - w_{N-1} - c_N)(v_1 - v_0) \\
&\Leftrightarrow r(w_1 - c_1) = v_1 \left( w_{N-1} - w_1 - \sum_{k=2}^{N-1} c_k + w_1 - c_1 + r - w_{N-1} - c_N \right) - \\
&\quad - v_0 \left( w_{N-1} - w_1 - \sum_{k=2}^{N-1} c_k + r - w_{N-1} - c_N \right) \\
&\Leftrightarrow v_1 = \frac{r(w_1 - c_1) + v_0 \left( r - w_1 - \sum_{k=2}^N c_k \right)}{r - \sum_{k=1}^N c_k} \tag{12}
\end{aligned}$$

$v_1$  in  $CR_1$  gives:

$$\begin{aligned}
CR_1 &= \frac{(w_1 - c_1) \left( r - \sum_{k=1}^N c_k \right)}{r(w_1 - c_1) + v_0 \left( r - w_1 - \sum_{k=2}^N c_k \right) - v_0 \left( r - \sum_{k=1}^N c_k \right)} = \\
&= \frac{(w_1 - c_1) \left( r - \sum_{k=1}^N c_k \right)}{r(w_1 - c_1) - v_0(w_1 - c_1)} = \frac{r - \sum_{k=1}^N c_k}{r - v_0} = CR_g \tag{13}
\end{aligned}$$

So  $CR_j = CR_g$  is true for every  $j$  between 1 and  $N$ . □

## 2 Horizontally interacting newsvendors

In order to transfer the vertical model to horizontal supply chain interaction we assume an assembly chain without substitution. Each production competence is unique in the supply chain. The parameters  $c_1, \dots, c_N$  and  $r$  are defined as in section 1. Site  $N$

assembles components, which are supplied by the sites 1 to N-1. The transfer prices  $w_1$  to  $w_{N-1}$  are paid by site N to the sites 1 to N-1 for one unit of a component to be assembled. The salvage value  $v_0$  is paid directly to the assembly site. The prices  $v_1$  to  $v_{N-1}$  represent transfer payments, which are introduced to share the overage risk. Vertical supply chain interaction between more than two value stages is not considered (cf. figure 2).

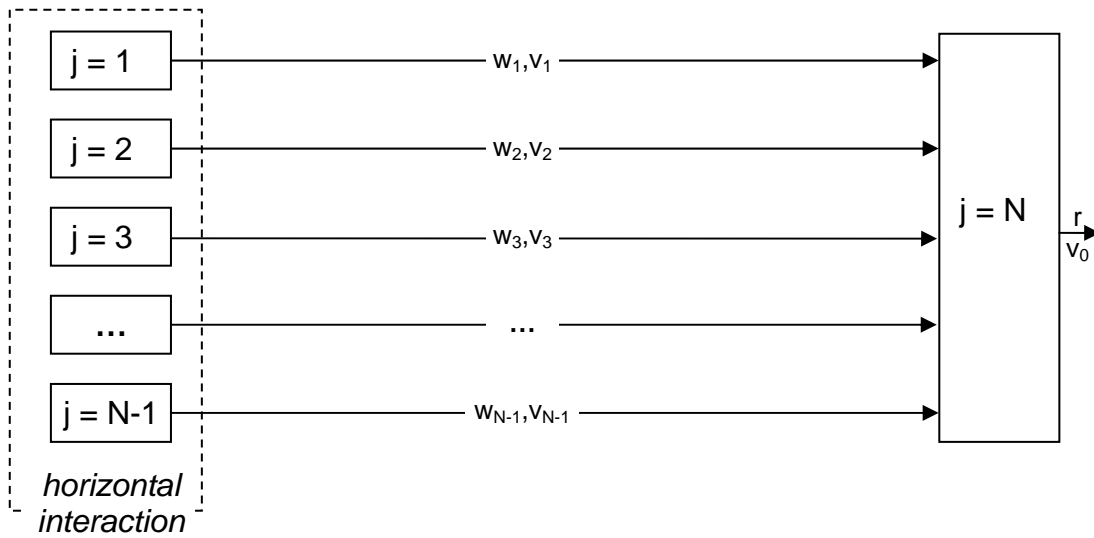


Figure 2: Decision situation for horizontal interaction

*Theorem 2:*

If  $w_1, \dots, w_{N-1}$  and  $v_1, \dots, v_{N-1}$  are determined in a way that  $CR_j = CR_{j-1}$  holds for every  $j$  between 2 and  $N$ , then  $CR_j = CR_0$  is true for every  $j$  between 1 and  $N$ , where

$$CR_j = \frac{w_j - c_j}{v_j} \quad \text{for } 1 \leq j \leq N-1 \quad (14)$$

and

$$CR_N = \frac{r - \sum_{k=1}^{N-1} w_k - c_N}{r - \sum_{k=0}^{N-1} v_k} \quad (15)$$

and

$$CR_g = \frac{r - \sum_{k=1}^N c_k}{r - v_0} \quad (16)$$

*Proof:*

First it is shown by induction that the statement

$$(CR_m = CR_{m-1} \quad \forall 2 \leq m \leq j) \Leftrightarrow v_j = \frac{w_j - c_j}{w_1 - c_1} v_1 \quad (17)$$

is true for every  $j$  between 2 and  $N-1$ .

Basis: Showing that the statement (17) holds for  $j = 2$ :

$$CR_2 = CR_1 \Leftrightarrow \frac{w_2 - c_2}{v_2} = \frac{w_1 - c_1}{v_1} \Leftrightarrow v_2 = \frac{w_2 - c_2}{w_1 - c_1} v_1 \quad (18)$$

Inductive step:

Induction hypothesis: Statement (17) holds for  $j$ .

Showing that Statement (17) then also holds for  $j+1$ :

$$\begin{aligned} CR_{j+1} &= CR_j \\ \Leftrightarrow \frac{w_{j+1} - c_{j+1}}{v_{j+1}} &= \frac{w_j - c_j}{v_j} \Leftrightarrow v_{j+1} = \frac{w_{j+1} - c_{j+1}}{w_j - c_j} v_j \end{aligned} \quad (19)$$

Using the induction hypothesis  $v_{j+1}$  can be rewritten to

$$\begin{aligned} v_{j+1} &= \frac{w_{j+1} - c_{j+1}}{w_j - c_j} v_j = \frac{(w_{j+1} - c_{j+1})(w_j - c_j)}{(w_j - c_j)(w_1 - c_1)} v_1 = \\ &= \frac{w_{j+1} - c_{j+1}}{w_1 - c_1} v_1 \end{aligned} \quad (20)$$

This way it is proven that statement (17) holds for every  $j$  between 2 and  $N-1$ . Now

we can calculate an expression for  $v_1$  from  $CR_N = CR_{N-1}$ :

$$\begin{aligned}
CR_N &= CR_{N-1} \\
\Leftrightarrow \frac{r - \sum_{k=1}^{N-1} w_k - c_N}{r - \sum_{k=0}^{N-1} v_k} &= \frac{w_{N-1} - c_{N-1}}{v_{N-1}} \tag{21} \\
\Leftrightarrow \left( r - \sum_{k=1}^{N-1} w_k - c_N \right) v_{N-1} &= (w_{N-1} - c_{N-1}) \left( r - \sum_{k=0}^{N-1} v_k \right)
\end{aligned}$$

Replacing all  $v_j$  with the expression from (17):

$$\begin{aligned}
\Leftrightarrow \left( r - \sum_{k=1}^{N-1} w_k - c_N \right) \frac{w_{N-1} - c_{N-1}}{w_1 - c_1} v_1 &= (w_{N-1} - c_{N-1}) \left( r - v_0 - \sum_{k=1}^{N-1} \frac{w_k - c_k}{w_1 - c_1} v_1 \right) \\
\Leftrightarrow \left( r - \sum_{k=1}^{N-1} w_k - c_N \right) v_1 &= (r - v_0)(w_1 - c_1) - \sum_{k=1}^{N-1} (w_k - c_k) v_1 \\
\Leftrightarrow v_1 \left[ r - \sum_{k=1}^{N-1} w_k - c_N + \sum_{k=1}^{N-1} (w_k - c_k) \right] &= (r - v_0)(w_1 - c_1) \tag{22} \\
\Leftrightarrow v_1 \left( r - \sum_{k=1}^N c_k \right) &= (r - v_0)(w_1 - c_1) \\
\Leftrightarrow v_1 &= \frac{(r - v_0)(w_1 - c_1)}{r - \sum_{k=1}^N c_k}
\end{aligned}$$

$v_1$  in  $CR_1$  gives:

$$\begin{aligned}
CR_1 &= \frac{(w_1 - c_1) \left( r - \sum_{k=1}^N c_k \right)}{(r - v_0)(w_1 - c_1)} = \\
&= \frac{r - \sum_{k=1}^N c_k}{r - v_0} = CR_g \tag{23}
\end{aligned}$$

So  $CR_j = CR_g$  is true for every  $j$  between 1 and  $N$ . □

### 3 Conclusion

In this paper we introduce the mathematical basis for a new and innovative solution method for the multi-location newsvendor problem. In particular, this method allows decentralized planning approaches in situations with incomplete information.